

- ...selective assignments (bad: $a[i] = 6$).
 - The goal of an imperative program is to **change the state** [of the machine].
 - The goal of a functional programs is to **evaluate** (reduce, simplify) **expressions**.
- ...in general, **updating** assignments ($y = x + 1$ good; $x = x + 1$ bad):
 - A variable in an imperative program: a name for a container.
 - There is no proper concept of “variable” in functional programs. What is called “variable” is a **name** for an expression.
- ...explicit pointers, storage management.
- ...input/output.
- ...control structures (loops, conditional statements).
- ...jumps (break, goto, exceptions).

- Expressions (**without** side effects).
 - Referential transparency (i.e., substitutivity, congruence).
- Definitions (of constants, functions).
 - Functions (almost as in mathematics).

	Math	Haskell
square		square

- Types (including higher-order, polymorphic, and recursively-defined types).
 - tuples, lists, and trees, shared sub-structures, implicit cycles.
- Automatic storage management (garbage collection).

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	Math	Haskell
square	$\mathbb{N} \rightarrow \mathbb{N}$	square :: Integer -> Integer
square(x)	$= x \times x$	square x = x * x

A function is defined by a set of **rewriting rules**.

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SESSIONS, SCRIPTS, EVALUATION

```
< godel:306/slides > ghci
GHCi, version 7.4.1: http://www.haskell.org/ghc/  :? for help
Loading package ghc-prim ... linking ... done.
Loading package integer-gmp ... linking ... done.
Loading package base ... linking ... done.
Prelude> 66
66
Prelude> 6 * 7
42
Prelude> square 35567
<interactive>:4:1: Not in scope:
'square'
Prelude> :load example
[1 of 1] Compiling Main
( example.hs, interpreted )
Ok, modules loaded: Main.
*Main> square 35567
1265011489
*Main> square (smaller (5, 78))
25
*Main> square (smaller (5*10, 5+10))
225
*Main>
```

THINK RECURSIVELY

Instructions for reading a book:

- **C:** "While not on the end cover repeat: read the current page, set the current page to the next page."
- **Functional:** "If on the end cover, stop. Otherwise, read the first page, then read **recursively** the rest of the book."
- Other examples:
 - To climb a ladder, step on the first rung and then climb (recursively) the rest of the ladder.
 - To eat a six-course meal, eat the first meal and then eat (recursively) the rest of the meal.
- How does one compute the factorial on a number?

WHAT'S LEFT? (CONT'D)

- Functions are **first order objects**.

```
twice :: (Integer -> Integer) -> (Integer -> Integer)
twice f = g
  where g x = f (f x)
```

- A program (or **script**) is a collection of definitions.
- Predefined data types in a nutshell:
 - Numerical: *Integer, Int, Float, Double*.
 - Logical: *Bool* (values: *True, False*).
 - Characters: *Char* ('a', 'b', etc.).
 - Composite: Functional: *Integer* \rightarrow *Integer*;
Tuples: *(Int, Int, Float)*;
Combinations: *(Int, Float)* \rightarrow *(Float, Bool)*, *Int* \rightarrow *(Int* \rightarrow *Int)*.

SCRIPTS

- Recall that a program is a collection of **definitions of values** (including functions).
- Syntactical sugar: definitions by **guarded equations**:

```
smaller :: (Integer, Integer) -> Integer
smaller (x,y)
  | x <= y  = x
  | x > y   = y
```

- Recursive definitions:

```
fact    :: Integer -> Integer
fact x = if x==0 then 1 else x * fact (x-1)
```

- Syntactical sugar: definitions by **pattern matching** (aka by cases):

```
fact    :: Integer -> Integer
fact 0 = 1
fact x = x * fact (x-1)
```

LOCAL DEFINITIONS

- Two forms:

```
let v1 = e1      def
  v2 = e2        where v1 = e1
  .             v2 = e2
  .             .
  .             .
  vk = ek        .
in exp          vk = ek
```

- Definitions are qualified by **where** clauses, while expressions are qualified by **let** clauses.

SCOPING

- Haskell uses **static** scoping.

```
cylinderArea :: Float -> Float -> Float
cylinderArea h r = h * 2 * pi * r + 2 * pi * r * r
```

```
cylinderArea1 :: Float -> Float -> Float
cylinderArea1 h r = x + 2 * y
  where x = h * circLength r
        y = circArea r
        circArea x = pi * x * x
        circLength x = 2 * pi * x
```

```
cylinderArea2 :: Float -> Float -> Float
cylinderArea2 h r = let x = h * circLength r
                    y = circArea r
                    in x + 2 * y
  where circArea x = pi * x * x
        circLength x = 2 * pi * x
```

TYPES

- Each type has associated operations that are not necessarily meaningful to other types.
 - Arithmetic operations (+, -, *, /) can be applied to numerical types, but it does not make any sense to apply them on, say, values of type *Bool*.
 - It does, however make sense to compare (using = (==), ≠ (/=), ≤ (<=), <, etc.) both numbers and boolean values.
- Every well formed expression can be assigned a type (**strong typing**).
 - the type of an expression can be inferred from the types of the constituents of that expression.
 - those expression whose type cannot be inferred are rejected by the compiler.

```
badType x      fact :: Integer -> Integer
| x == 0 = 0    fact x
| x > 0 = 'p'   | x < 0 = error "Negative argument."
| x < 0 = 'n'   | x == 0 = 1
                | x > 0 = x * fact (x-1)
```

What is the type of error?

TWO DATA TYPES

- Booleans.** Values: *True*, *False*.
 - operations on *Bool*: logic operators: ∨ (||), ∧ (&&), ¬ (not); comparisons: = (==), ≠ (/=); relational <, ≤ (<=), >, ≥ (>=).
- Characters.** Values: 256 of them, e.g., 'a', 'b', '\n'.
 - Operations on characters: comparison, relational;

```
ord :: Char -> Int      Prelude> import Data.Char
chr :: Int -> Char      Prelude Data.Char> ord 'a'
                        97
                        Prelude Data.Char> chr 100
                        'd'

toLower :: Char -> Char
toLower c | isUpper c = chr (ord c - (ord 'A' - ord 'a'))
          | True      = c
          where isUpper c = 'A' <= c && c <= 'Z'
```

LISTS

- A list is an ordered set of values.

<code>[1, 2, 3] :: [Int]</code>	<code>[[1, 2], [3]] :: [[Int]]</code>	<code>['h','i'] :: [Char]</code>
<code>[div, rem] :: ??</code>	<code>[1,'h'] :: ??</code>	<code>[] :: ??</code>

- Syntactical sugar:

```
Prelude> ['h','i']
"hi"
Prelude> "hi" == ['h','i']
True
Prelude> [['h','i'], "there"]
["hi", "there"]
```

CONSTRUCTING LISTS

- Constructors: `[]` (the empty list) and `:` (constructs a longer list).

```
Prelude> 1:[2,3,4]
[1,2,3,4]
Prelude> 'h': 'i':[]
"hi"
```

- The operator `:` (pronounced “cons”) is **right associative**.
- The operator `:` **does not** concatenate lists together!

```
Prelude> [1,2,3] : [4,5]
No instance for (Num [t0])
arising from the literal `4'
Possible fix: add an instance declaration for (Num [t0])
In the expression: 4
In the second argument of `(:)', namely `[4, 5]'
In the expression: [1, 2, 3] : [4, 5]
Prelude> [1,2,3] : [[4,5]]
[[1,2,3],[4,5]]
Prelude> [1,2,3] ++ [4,5]
[1,2,3,4,5]
Prelude>
```

OPERATIONS AND PATTERN MATCHING ON LISTS

- Comparisons (`<`, `>`, `==`, etc.), if possible, are made in lexicographical order.
- Subscript operator: `!!` (e.g., `[1, 2, 3] !! 1` evaluates to 2) – **expensive**
- Arguably the most common list processing: Given a list, do something with each and every element of that list.
 - In fact, such a processing is so common that there exists the predefined *map* that does precisely this:

```
map f [] = []
map f (x:xs) = f x : map f xs
```
 - This is also an example of **pattern matching on lists**.
 - * Variant to pattern matching: *head* and *tail* (predefined).

```
head (x:xs) = x      | map f l = if l == []
tail (x:xs) = xs     |      then []
                      |      else f (head l) : map f (tail l)
```

TUPLES

- While lists are homogenous, tuples group values of (possibly) different types.

```
divRem :: Integer -> Integer -> (Integer, Integer)
divRem x y = (div x y, rem x y)
```

```
divRem1 :: (Integer, Integer) -> (Integer, Integer)
divRem1 (x, 0) = (0, 0)
divRem1 (x, y) = (div x y, rem x y)
```

- The latter variant is also an example of pattern matching on tuples.

- An operator contains symbols from the set `!#$%&*+./<=>?@\`|`: (– and ~ may also appear, but only as the first character).
- Some operators are predefined (+, –, etc.), but you can define your own as well.
- An (infix) operator becomes (prefix) function if surrounded by brackets. A (prefix) function becomes operator if surrounded by backquotes:

```
divRem :: Integer -> Integer -> (Integer, Integer)
x `divRem` y = (div x y, rem x y)
-- precisely equivalent to
-- divRem x y = (div x y, rem x y)

(%%) :: Integer -> Integer -> (Integer, Integer)
(%%) x y = (div x y, rem x y)
-- precisely equivalent to
-- x %% y = (div x y, rem x y)
```

These are just lexical conventions.

- Identifiers consist in letters, numbers, simple quotes ('), and underscores (_), but they **must** start with a letter.
 - For the time being, they must actually start with a **lower case letter**.
 - * A Haskell identifier starting with a capital letter is considered a type (e.g., *Bool*) or a type constructor (e.g., *True*)—we shall talk at length about those later.
 - * By convention, types (i.e., class names) in Java start with capital letters, and functions (i.e., method names) start with a lower case letter. What is a convention in Java is **the rule** in Haskell!
 - Some identifiers are language keywords and cannot be redefined (if, then, else, let, where, etc.).
 - * Some identifiers (e.g., *either*) are defined in the standard prelude and possibly cannot be redefined (depending on implementation, messages like “Definition of variable “either” clashes with import”).