

# CS 316: Constraint satisfaction problems

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- Standard search problem: the state is anything that supports goal test, comparison, successor
- CSP: the state is defined by **variables**  $V_i$  with values from **domains**  $D_i$ 
  - The **goal test** is a set of **constraints**, which specifies allowable combinations of values for subsets of variables. A state is a set of variable bindings.

*Variables*    {    *Shoes, Pants, Shirt*    }

*Domains*    {    {*Sandals, Runners*}, {*Jeans, Blue, Grey*}, {*Green, White*}    }

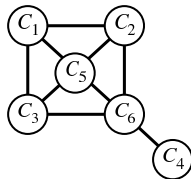
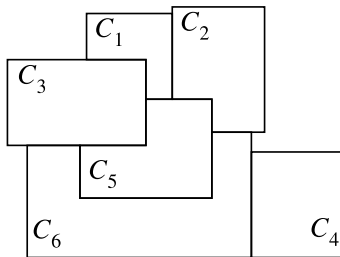
*Constraints*    {    (*Shoes = Sandals, Pants = Grey*),  
                          (*Shoes = Runners, Pants = Jeans*),  
                          (*Shoes = Sandals, Shirt = Green*),  
                          (*Shoes = Runners, Shirt = White*),  
                          (*Pants = Grey, Shirt = Green*),  
                          (*Pants = Jeans, Shirt = White*),  
                          (*Pants = Blue, Shirt = White*)    }

- This is actually an example of **binary CSP**

# EXAMPLE: MAP COLOURING



- Colour a map so that no adjacent countries have the same colour.



- Variables: Countries  $C_i$
- Domains:  $\{Red, Green, Blue\}$
- Constraints:  $C_1 \neq C_2$ ,  $C_1 \neq C_3$ ,  $C_3 \neq C_5$ , ...



- Assignment problems
  - e.g., who teaches what class
- Timetabling problems
  - e.g., which class is offered when and where?
- Hardware configuration
- Spreadsheets
- Transportation scheduling
- Factory scheduling

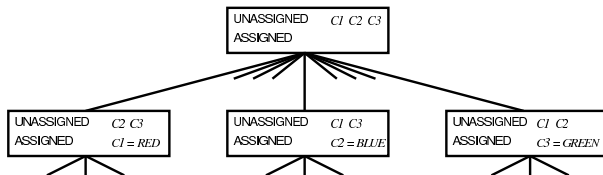


- States are defined by the variables bound so far.

**Initial state** All variables unbound.

**Operators** Bind one variable

**Goal test** All variables assigned, no constraints violated



- Disadvantages?



- Order of assignment is irrelevant (many paths are equivalent)
- Further bindings cannot correct an already violated constraint
- We can use depth-first search, but
  - Fix the order of assignment
  - Check for constraint violations
    - at the SUCCESSORS level, or immediately before expanding the state.

```
function CSP-SEARCH() returns a solution, or failure  
  
nodes ← MAKE-QUEUE(MAKE-NODE(INITIAL-STATE))  
repeat  
  if nodes is empty then return failure  
  node ← REMOVE-FRONT(nodes)  
  if GOAL-TEST(node) then return node  
  unless VIOLATES-CONSTRAINTS(node) do  
    nodes ← APPEND(SUCCESSORS(node), nodes)  
forever
```

- We do not need the queue actually (why?)

- We fix the order of assignment, and we check for constraint violations
- The resulting algorithm is called **backtracking**, the basic uninformed algorithm for CSP. Can solve  $n$ -queens for  $n \approx 15$

**function** BACKTRACKING(*state, variables, domains*) **returns** a solution, or failure

*var*  $\leftarrow$  FIRST(*variables*)

*domain*  $\leftarrow$  FIRST(*domains*)

**foreach** *val* **in** *domain* **do**

**unless** VIOLATES-CONSTRAINTS(ADD((*var, val*), *state*)) **do**

**if** no more variables **then return** *state*

        BACKTRACKING(ADD((*var, val*), *state*),  
                        REST(*variables*), REST(*domains*))

**end**

**return** failure

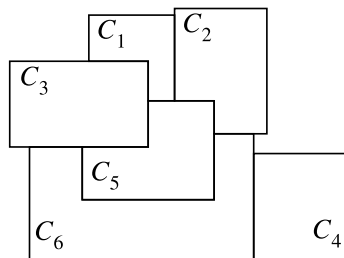


- Variant: **forward checking**, looks ahead and erases from the domains of all the variables those values that cannot be assigned without violating constraints
  - Forward checking is a particular case of **arc consistency** working on the graph generated by the constraints
  - Arc consistency can be applied on the initial graph of constraints before performing the backtracking search



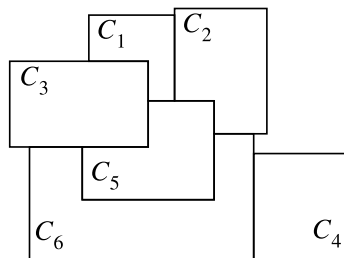


- We can make more intelligent decisions on
  - which value to choose for each variable
  - which variable to assign next
- Given  $C_1 = \text{Red}$  and  $C_2 = \text{Green}$ ,  $C_3 = ?$
- Given  $C_1 = \text{Red}$ ,  $C_2 = \text{Green}$ , what next?





- We can make more intelligent decisions on
  - which value to choose for each variable
  - which variable to assign next
- Given  $C_1 = \text{Red}$  and  $C_2 = \text{Green}$ ,  $C_3 = ?$ 
  - $C_3 = \text{Green}$ , the least constraining value
- Given  $C_1 = \text{Red}$ ,  $C_2 = \text{Green}$ , what next?
  - Choose  $C_5$ , the most constrained variable



- Can solve  $n$ -queens for  $n \approx 1000$