

CS 406: Compilers and Interpreters

Stefan D. Bruda

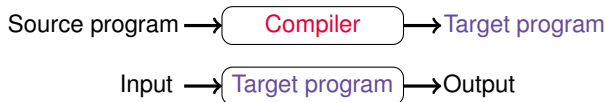
Winter 2016



- Coordinates:
 - **Course Web page:**
<http://cs.ubishops.ca/home/cs406>
 - Instructor: Stefan Bruda
(<http://bruda.ca>, stefan@bruda.ca, Johnson 114B, ext. 2374)
 - Office hours?
- Textbook (required): C. N. Fischer, R. K. Cytron, and R. J. LeBlanc Jr, *Crafting a Compiler*, Addison Wesley, 2009.



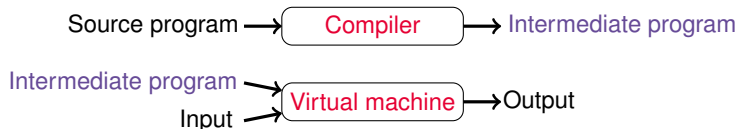
- **Pure compilation:** The compiler translates the high-level source program into an equivalent target program (typically in machine language), then goes away:



- **Pure interpretation:** The interpreter stays around for the execution of the program and becomes the locus of control during execution



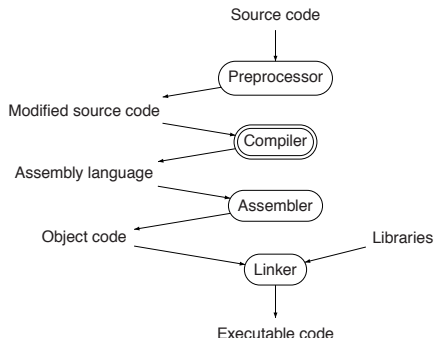
- **Compilation followed by interpretation:**



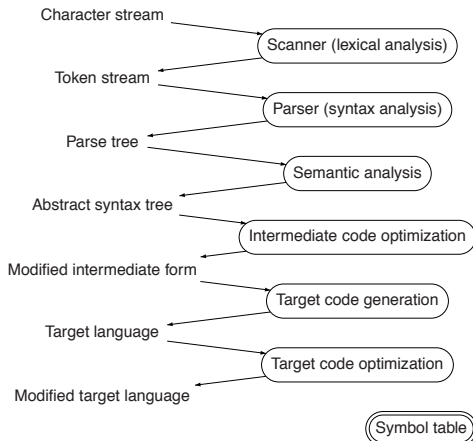


- Interpretation offers greater flexibility and better diagnostics, but compilation offers better performance
- Compilation does **not** have to produce machine language for some hardware
 - **Compilation = translation from one language into another**
 - Some compilers produce nothing but virtual instructions (Pascal P-code, Java byte code, Microsoft COM+)
- Compilation possibly preceded by a **preprocessor**

- For languages that compile to executable code:



- For languages that run on a virtual machine: the assembler and linker part are replaced by an interpreter (or virtual machine)



- **Scanner**: divides program into “tokens” (smallest meaningful units)
 - Driven by **regular expressions**
- **Parser**: discovers the syntactic structure of a program
 - Driven by **context-free grammar**
- **Semantic analysis**: discovers the meaning of the program
 - **Static analysis**
 - Some other things can only be figured out at run time
- **Intermediate form**: tree-like structure and/or some machine-like language (but machine independent)
 - Often a form of machine language, but for an idealized machine



- **Intermediate code optimization**: produce code that does the same thing, only faster
 - Algorithmic optimization
- **Code generation**: produces assembly language for the target machine
- **Code optimization**: machine-specific optimizations (use of special instructions or addressing modes, reorder instruction to improve the load on superscalar architectures, etc.)
- **Symbol table**: all phases rely on a symbol table that keeps track of all the identifiers in the program and what the compiler knows about them
 - This symbol table may be retained (in some form) even after compilation has completed, for use by a debugger