Concurrency

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- Concurrency can be achieved by multiprocessing and time-sharing
 - Best definition for concurrency: apparently simultaneous execution
- Concurrency is fundamental to distributed computing
 - Multiprocessing: many machines run simultaneously many programs
 - Time-sharing: a single machine runs multiple programs in an interleaved fashion (context switching)
- Whether things run in a time-share fashion or on a multiprocessor computer is immaterial; the observable behaviour is the same



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- Whether things run in a time-share fashion or on a multiprocessor computer is immaterial; the observable behaviour is the same
- The fundamental unit of computation: a process
 - A process consists of an address space and one (or more) threads of execution (i.e., instruction pointers)
 - Each process receives a separate copy of all the variables (address space)
 - Each thread in a process have a copy of local variables (stack) but they all share the rest of the address space (global variables, heap)



```
• We now create singly threaded processes:
```

#include <iostream.h>
#include <sys/types.h>
#include <unistd.h>

```
int main (int argc, char** argv) {
    int i;
    int sum = 0;
    fork();
    for (int i=1; i<10000; i++) {
        sum = sum + i;
     }
        cout << "\nI computed " << sum << "\n";
}</pre>
```

to To other side. get the Why did the multithreaded chicken cross the road? other to side. To the get



- The call to fork() duplicates the current process; both processes continue execution from the instruction following the call to fork()
- The initial process is the parent, and the newly created copy is called a child
- Processes are identified in a Unix system by a unique process identifier (or PID, unsigned integer)
- fork() returns two different integers in the child and parent processes
 - In the child process fork() returns zero
 - In the parent process fork() returns the PID of the newly created child
 - So we can provide different code for the parent and the child, by surrounding them in appropriate conditional statements

DIVERGING PROCESSES

```
#include <iostream.h>
#include <sys/types.h>
#include <unistd.h>
int main (int argc, char** argv) {
  int i; int sum = 0; int pid;
  pid = fork();
  for (int i=1; i<10000; i++) {
    if (pid == 0)
      cout << "+"; // child process</pre>
    else
      cout << "-"; // parent process</pre>
    cout.flush();
    sum = sum + i; // both processes
  }
  if (pid == 0)
    cout << "\n[Child] I computed " << sum << "\n";</pre>
  else
    cout << "\n[Parent] I computed " << sum << "\n";</pre>
}
```



- The call execve replaces completely the current process with another executable
 - The arguments are the name of the command to execute, then two null-terminated arrays of strings containing the command line arguments and the environment, similar to the ones received by the function main
- Suppose now that we want to run an external command (so we use execve), but also we want to continue the execution of the original program



- The call execve replaces completely the current process with another executable
 - The arguments are the name of the command to execute, then two null-terminated arrays of strings containing the command line arguments and the environment, similar to the ones received by the function main
- Suppose now that we want to run an external command (so we use execve), but also we want to continue the execution of the original program
- We use a combination of fork and execve:

```
int childp = fork();
if (childp == 0) { // child
  execve(command, argv, envp);
}
else { // parent
  // code that continues our program
}
```



- Again, suppose that we want to execute an external command (execve again)
- We still want to continue the execution of the main program
- But only after the run of the external command is complete (synchronous as opposed to asynchronous execution)



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- We still want to continue the execution of the main program
- But only after the run of the external command is complete (synchronous as opposed to asynchronous execution)

```
int run_it (char* command, char* argv [], char* envp[]) {
  int childp = fork();
  int status;
```

```
if (childp == 0) { // child
  execve(command, argv, envp);
}
else { // parent
  waitpid(childp, &status,0);
}
return status;
```

}

- We will build eventually servers (programs that serve requests from clients)
- Instead of serving requests from one client at a time, our servers will handle many clients quasi-simultaneously

```
loop
| listen for clients
| if a client requests connection then
| fork
| if child process then
| handle client
terminate
forever
```

- We will build eventually servers (programs that serve requests from clients)
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```
loop

listen for clients

if a client requests connection then

fork

if child process then

handle client

terminate

forever
```

... or even many types of clients

```
loop

listen for clients

if a client of type x requests connection then

fork

if child process then

launch server x

terminate

forever
```





